Integral Cryptanalysis Using Algebraic Transition Matrices

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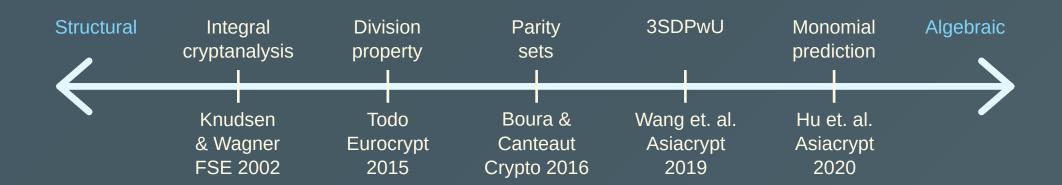
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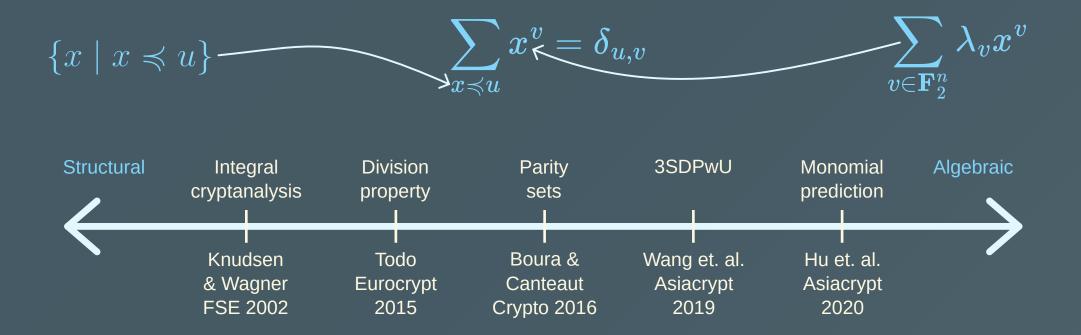
What's in a title: Integral Cryptanalysis



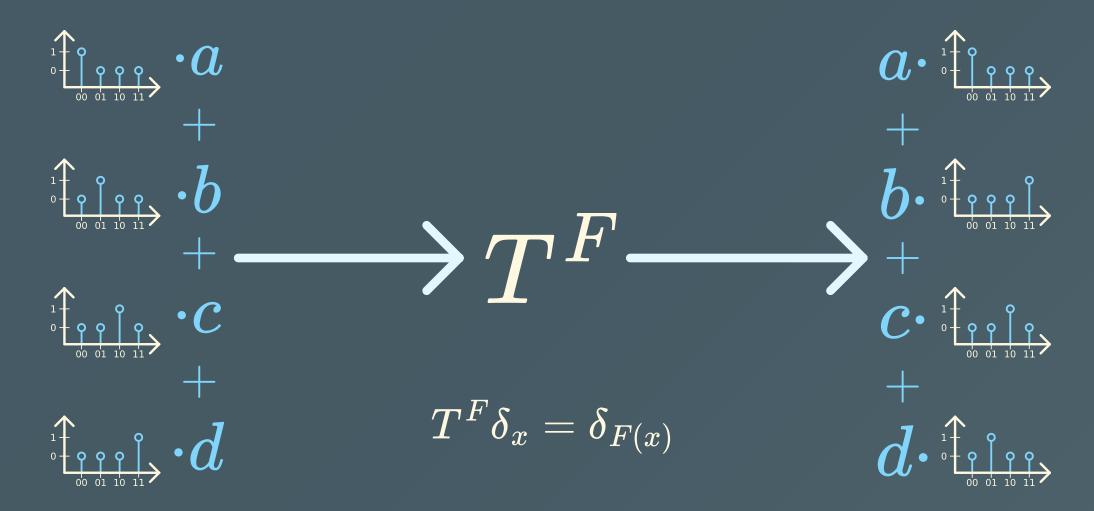
A Brief History of Integral Cryptanalysis



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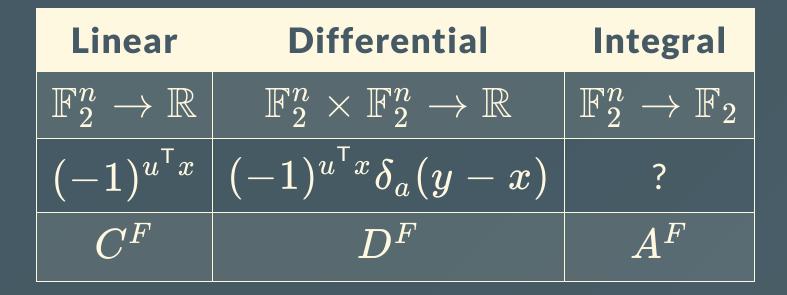
What's in a title: Transition Matrices

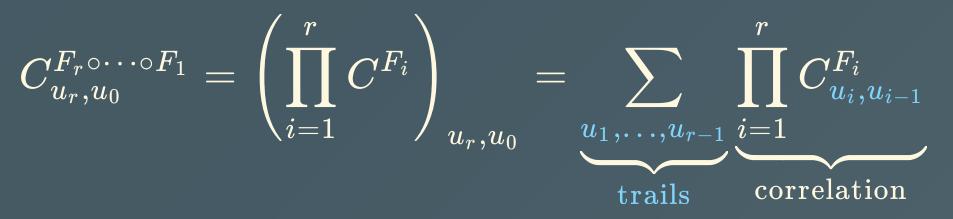


What's in a title: Transition Matrices

Linear	Differential	Integral
$\mathbb{F}_2^n o \mathbb{R}$	$\mathbb{F}_2^n imes \mathbb{F}_2^n o \mathbb{R}$	$\mathbb{F}_2^n o \mathbb{F}_2$
$(-1)^{u^{T}x}$	$(-1)^{u^{ op}x}\delta_a(y-x)$?
C^F	D^F	A^F

What's in a title: Transition Matrices





Algebraic Transition Matrices

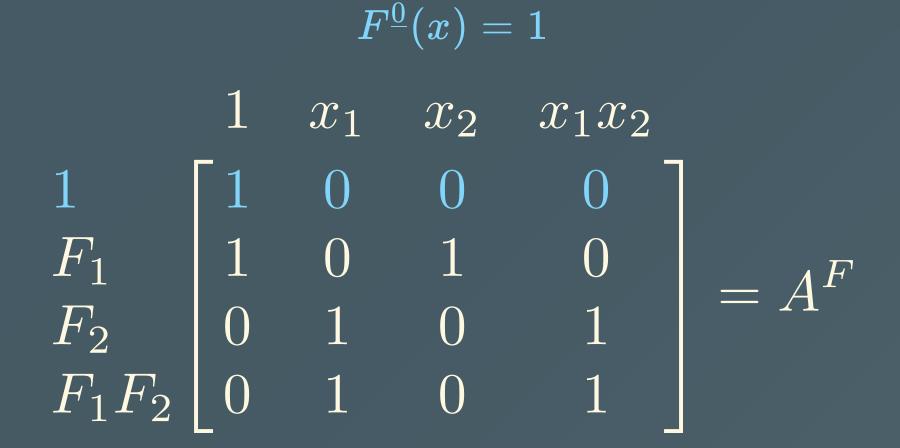
$$\sum_{x \preccurlyeq u} x^v = \underbrace{\mathbf{1}_{\{x \preccurlyeq u\}}}_{\text{precursor}} \cdot \underbrace{x^v}_{\text{monomial}} = \delta_{u,v}$$

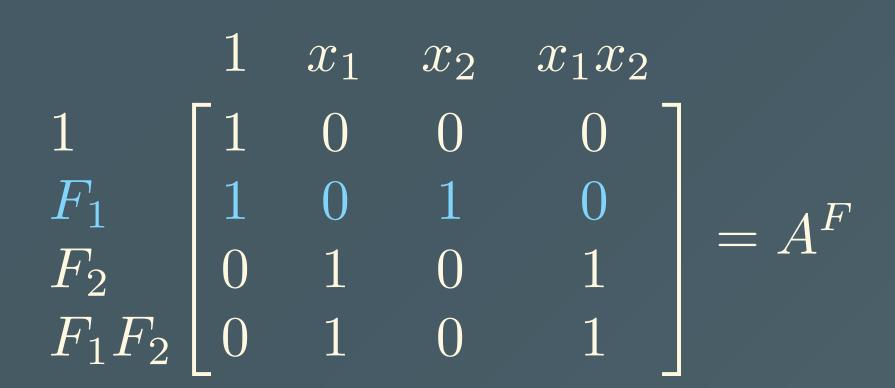
$$A^F = \mathscr{P}_m T^F \mathscr{P}_n^{-1}$$

$${\mathscr P}_n={\mathscr P}_n^{-1}=egin{bmatrix} 1&1\0&1\end{bmatrix}^{\otimes r}$$

 $F(x)=(F_1(x),F_2(x))=(x_2+1,x_1+x_1x_2)$







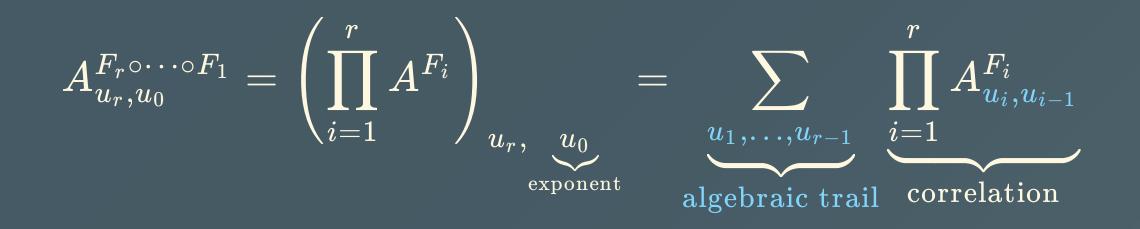
 $F_1(x) = x_2 + 1$

 $F_2(x) = x_1 + x_1 x_2$ $1 \quad x_1 \quad x_2 \quad x_1 x_2$ $\begin{bmatrix} 1 & 0 & 0 & 0 \\ F_1 & 0 & 1 & 0 \\ F_2 & 0 & 1 & 0 & 1 \\ F_1 F_2 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 \end{bmatrix} = A^F$

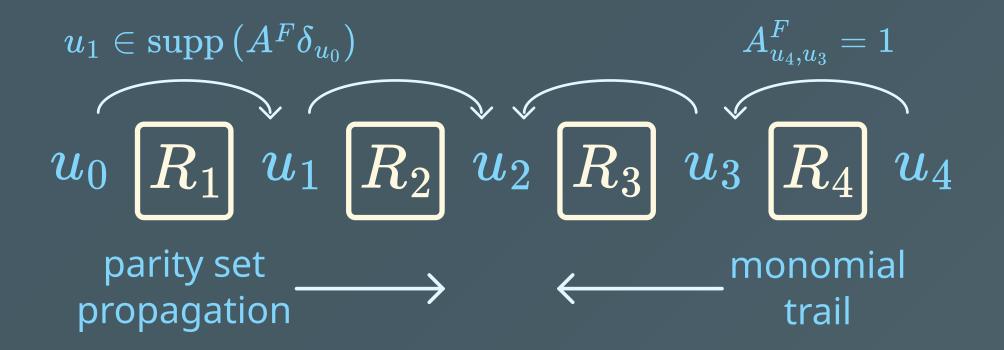
 $1 \quad x_1 \quad x_2 \quad x_1 x_2$ $\begin{bmatrix}
1 & 0 & 0 & 0 \\
F_1 & 0 & 1 & 0 \\
F_2 & 0 & 1 & 0 & 1 \\
F_1 F_2 & 0 & 1 & 0 & 1 \\
0 & 1 & 0 & 1
\end{bmatrix} = A^F$

 $F^{1}(x) = x_1 + x_1 x_2$

Algebraic Trails



Algebraic Trails



Key Addition

$$egin{aligned} &A^{ au_k} = \bigotimes_i egin{bmatrix} 1 & 0 \ k_i & 1 \end{bmatrix} \ &A^{ au_k}_{v,u} = egin{cases} &k^{u+v} & ext{if } u \preccurlyeq v \ 0 & ext{otherwise} \end{aligned}$$

 $A_{v,u}^{ au_{k_2}\circ F\circ au_{k_1}}
eq 0 \implies A_{v'\succcurlyeq v,u'\preccurlyeq u}^{ au_{k_2}\circ F\circ au_{k_1}}=f(k_1,k_2)$

Finding Integral Properties Simple properties

$$\sum_{x \preccurlyeq u} F^v(x) = A^F_{v,u}$$

1. Prove all trails from u to v have zero correlation.

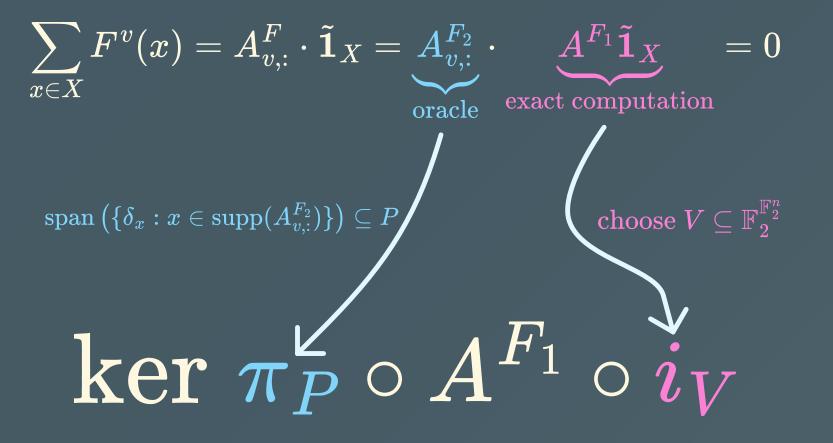
2. Prove all trails from u to v are key-independent, then sum trails.

3. Sum trails from u to v per key-monomial.

$$\sum_{x\in X}r\left(F(x)
ight)=\left({\mathscr P}_m^{-{\sf T}}r
ight)\cdot A^F\left({\mathscr P}_n{f 1}_X
ight)=0,$$

$$\sum_{x\in X}r\left(F(x)
ight)=r^{\circ}\cdot A^{F} ilde{\mathbf{1}}_{X}=0 \ =\operatorname{vec}\left(A^{F}
ight)\cdot\left(ilde{\mathbf{1}}_{x}\otimes r^{\circ}
ight) \ =\operatorname{vec}\left(A^{F_{k_{1}}}
ight) \ = 0 \ \stackrel{f^{k_{1}}}{\underset{k_{n}}{\overset{k_{n}}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_{n}}{\overset{k_$$

$$\sum_{x\in X}F^v(x)=A^F_{v,:}\cdot ilde{\mathbf{1}}_X=A^{F_2}_{v,:}\cdot \quad A^{F_1} ilde{\mathbf{1}}_X=0$$



Finding Integral Properties on 9-round PRESENT

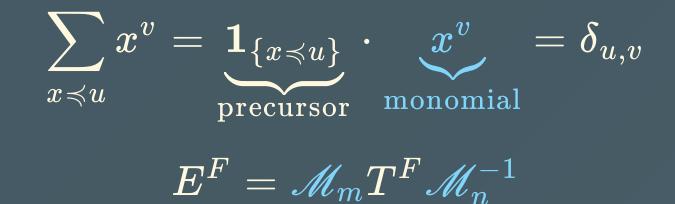
• 9 rounds: 470 dimensional space of properties

$$\sum_{x_1=x_2=x_3=x_4=0}F_5(x)+F_{13}(x)=c_1
onumber \ \sum F_5(x)+\sum F_{12}(x)=c_2$$

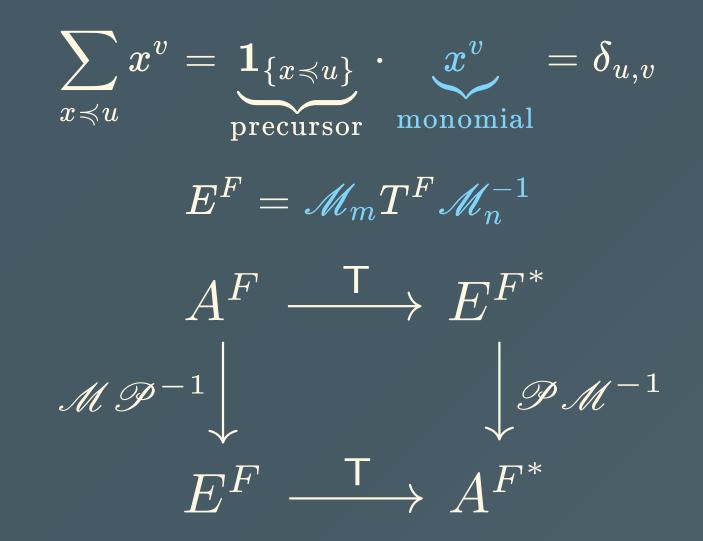
$$\sum_{x_5=0}F_5(x)+\sum_{x_9=0}F_{13}(x)=c_2$$

 $\sum_{x_5+x_9=0}F_1(x)F_{17}(x)F_{33}(x)F_{49}(x)=c_3$

Duality



Duality



Conclusion

- Integral cryptanalysis fits in the geometric approach
 - New insight in and better understanding of integral cryptanalysis
 - Improved search methods

Future work

- Don't ignore the key
 - Weak key
 - Key schedule
- Build on/Improve search for generalized integral properties
 - $\circ\,$ Allow key in computation
 - Key-recovery by selection of useful properties from solution space

Coming Soon

- Ultrametric integral cryptanalysis
 - $\circ\,$ Justification of basis by simplification of multiplication

$$\circ \sum_{x\in X} r(F(x)) = 0 \mod 2^l$$